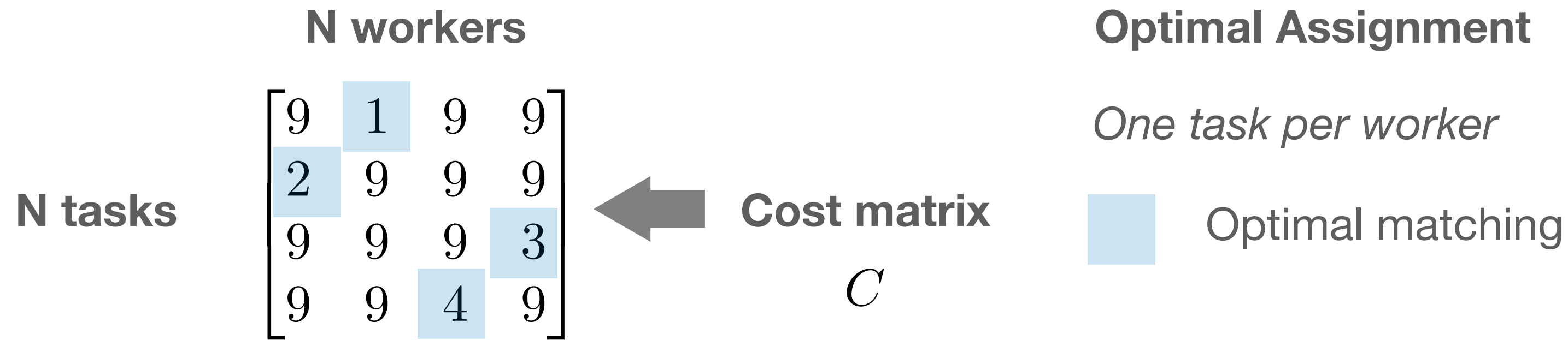
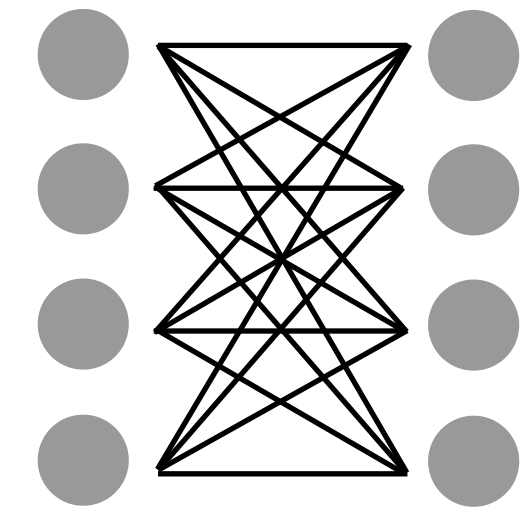


Hungarian (Munkres) Algorithm - Convex Relaxation



Bipartite Matching Perspective

N tasks N workers



...Costs are edge weights

If a worker can't perform a task...
either remove the edge...
or make the cost infity

Overall Cost: Sum of costs for each assigned worker $\text{Tr}(C^T P)$ **do you see why?**

for some permutation matrix: $P \in \mathcal{P}$ ← not a convex set **do you see why?**

A permutation matrix is an identity matrix with the columns (or rows) shuffled... ex. $\begin{bmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$

Convex Relaxation: $\mathcal{P} \rightarrow \mathcal{X}$ Doubly stochastic matrices: rows and columns sum to 1 \mathcal{X} is convex ✓

Birkhoff-von Neumann Theorem: "the convex hull of permutation matrices is doubly stochastic matrices" **do you see why?**

Doubly stochastic matrices: $X \in \mathcal{X}$ satisfies $X\mathbf{1} = \mathbf{1}$ $X^T\mathbf{1} = \mathbf{1}$ $X \geq 0$

Convex Relaxation

$$\begin{aligned} \min_X & \text{Tr}(C^T X) \\ \text{s.t.} & X\mathbf{1} = \mathbf{1} \quad X \geq 0 \\ & X^T\mathbf{1} = \mathbf{1} \end{aligned}$$

Dual Program

$$\begin{aligned} \max_{v,w,U} & v^T\mathbf{1} + w^T\mathbf{1} \\ \text{s.t.} & C^T = \mathbf{1}v^T + w\mathbf{1}^T + U^T \\ & U \geq 0 \end{aligned}$$

Dual Variable Interpretation

- v_i : row potential - lower bound on cost of performing task i
- w_j : column potential - lower bound on cost incurred by worker j
- U_{ij} : Inefficiency of worker j performing task i

Note: this convex relaxation works cause linear programs are always optimized at the vertices

do you see why?